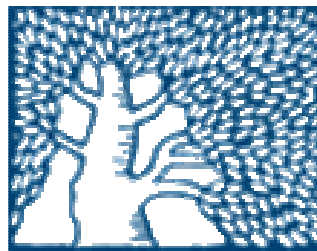


An Optimally Fair Coin Toss

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Coin Flipping

- Mutually distrustful parties want to flip a fair (binary) coin
- One party may be malicious
 - Can arbitrarily deviate from the protocol
 - In particular -- can **abort prematurely**
- Output of the honest party should not be significantly biased



Coin Flipping

- When the parties are **computationally unbounded**, one of them can control the binary coin

Without simultaneity

Lots of work: with more than two parties and **honest majority**

- Coin flipping implies **one-way functions**

[Impagliazzo-Luby '89]

- Blum '81: Coin flipping using **bit commitment**

This Work

- **Cleve '86:**
Any r -round protocol is $\Omega(1/r)$ biased
- **Best previously known protocol: bias $O(1/\sqrt{r})$**
 - Cleve-Impagliazzo '93: An $\Omega(1/\sqrt{r})$ lower bound in a tightly related model

Our result:

Cleve's bound is tight!

- **Construct an optimally fair protocol - bias $O(1/r)$**
 - Based on standard cryptographic assumptions (Oblivious Transfer)
 - Builds upon recent progress in **fair secure computation**
[Gordon-Hazay-Katz-Lindell '08]

Exact constants: between
 $1/4$ and $1/8$

The Main Idea

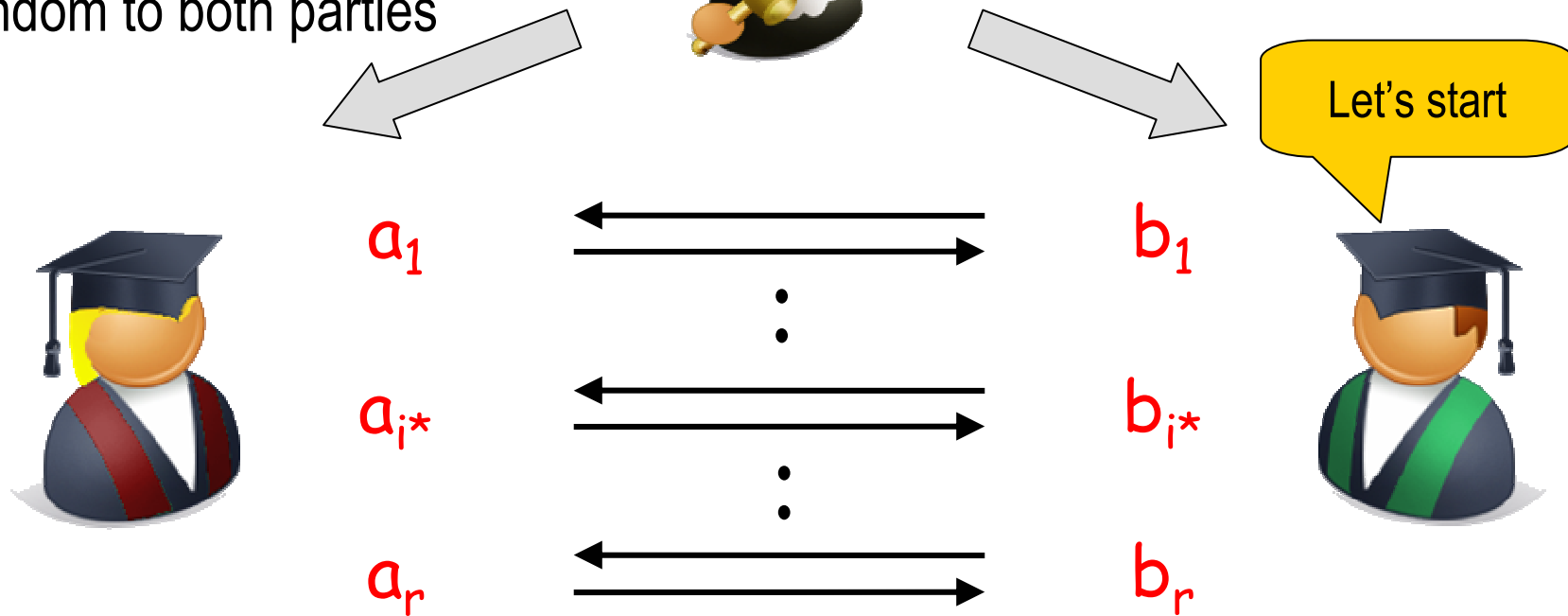
Threshold round

Add a new character: the dealer

At this point: coin should still be random to both parties



1. Chooses a secret $i^* \in [r]$
2. Distributes shares



- Until round i^* the bits a_i and b_i are random & independent
- $a_{i^*} = b_{i^*}$ is the output ($a_{i^*+1} = b_{i^*+1} = \text{"Halt"}$)
- If Alice halts at round $i \leq i^*+1$, Bob outputs b_{i-1}

Open Problems

- **Minimal assumptions for achieving the optimal $O(1/r)$ bias**

- Blum's protocol relies on **any** one-way function
- Our protocol relies on Oblivious Transfer

They are black-box separable

- **Efficient implementations**

- Pre-processing phase: relies on **general secure computation**
- The dealer's functionality is rather simple
- Is there a **specific** and more efficient implementation?

- **Optimal bias in the multiparty setting**

- Several straightforward extensions of our protocol fail
 - The adversary may increase the probability of guessing the crucial round
 - Can get bias $O(k/\sqrt{r})$

- **Fair protocols for other functionalities - GK2008**

Advertisements

- Games For Extracting Randomness
- <http://math166-pc.weizmann.ac.il/>
- <http://www.wisdom.weizmann.ac.il/~neko/>
 - Play a game to help science
- Weizmann Winter School on Foundations of Computer Science: **February 15-19th 2009**
 - **Tentative!**

תודה רבה
Thank you